## pTunes: Runtime Parameter Adaptation for Low-power MAC Protocols

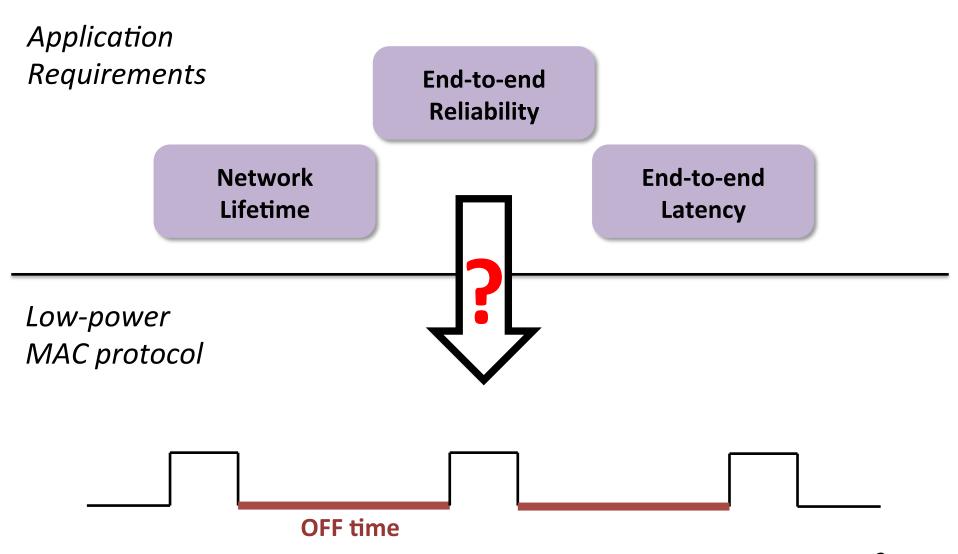
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#### Configuring a MAC Protocol is Not Easy



#### A Real-world Example

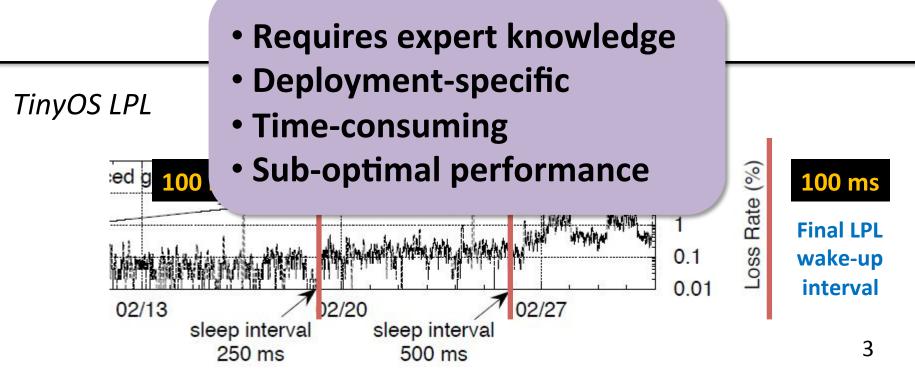
Adaptive Control Application

Ceriotti et al., IPSN'11

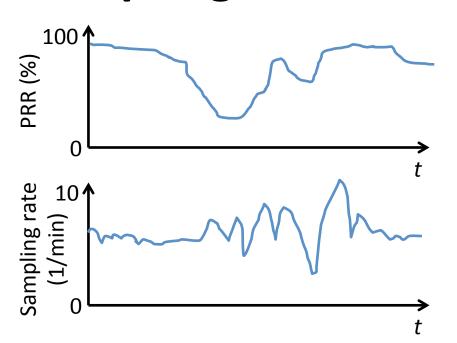


#### **Current practice:**

- Experience
- Field trials
- Over-provision



#### Adapting a MAC Protocol is Even Harder

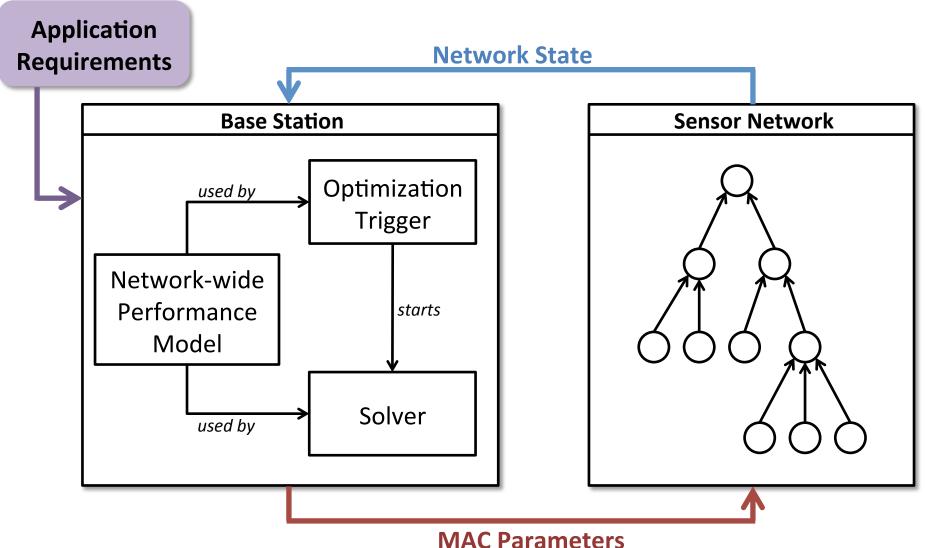




#### Need to adapt at runtime to changes in

- Topology (e.g., node failures, disconnects)
- Wireless link quality (e.g., interference)
- Traffic load (e.g., varying sampling rate)

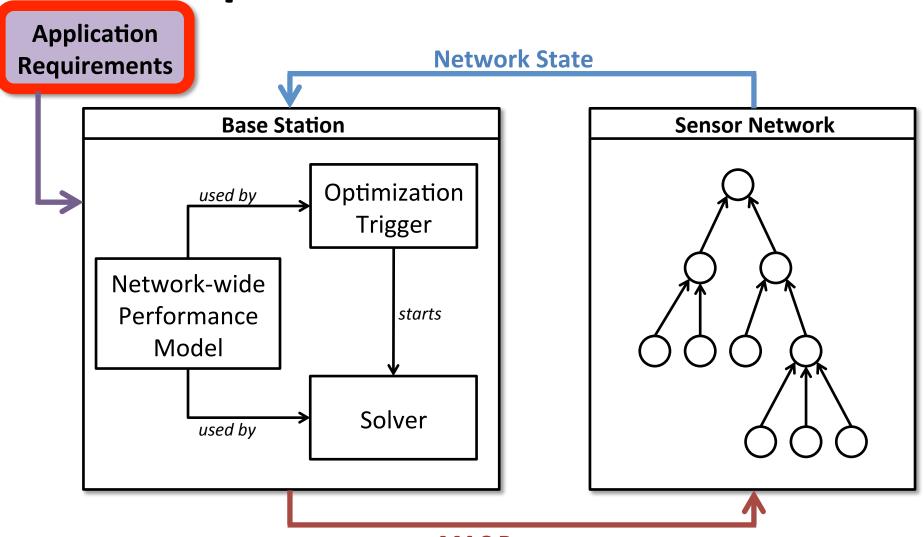
## pTunes in a Nutshell



#### **Contributions**

- 1. pTunes framework for runtime adaptability of existing low-power MAC protocols
- 2. Flexible modeling approach
- 3. Efficient runtime support to "close the loop"

## pTunes in a Nutshell



#### **Application Requirements**

- pTunes targets data collection scenarios
  - Tree routing
  - Low-power MAC
- Example requirements specification

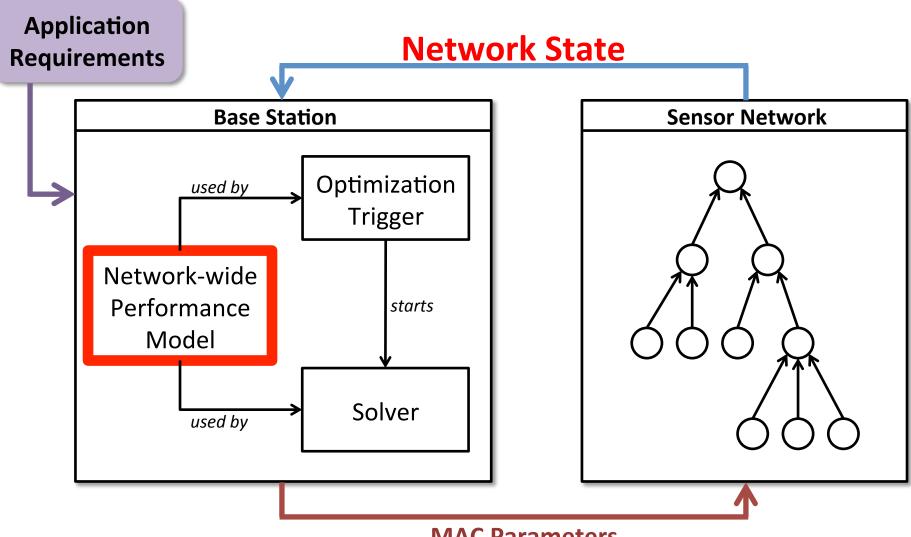
Maximize: Network lifetime

Subject to: End-to-end reliability greater than 95 %

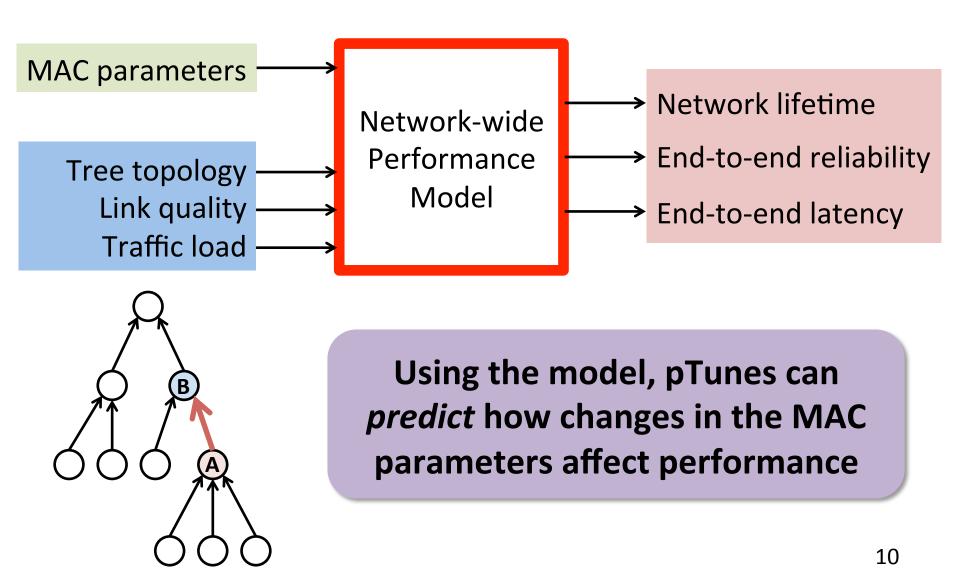
End-to-end latency below 1 second

pTunes determines at runtime MAC parameters whose performance meets the requirements

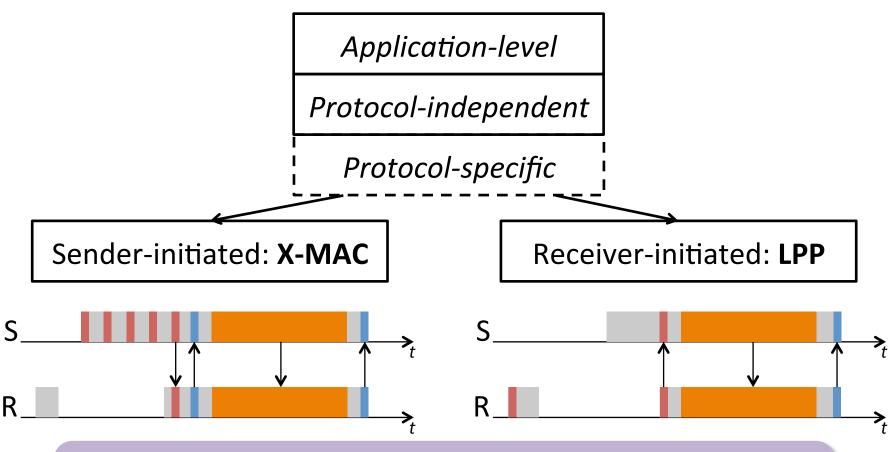
## pTunes in a Nutshell



#### **Network-wide Performance Model**



## **Layered Modeling Approach**



Only the lowest layer must be changed to adapt the model in pTunes to a given MAC protocol

# Layering in Action: X-MAC End-to-end Reliability

$$R = \frac{1}{|\mathcal{M}|} \sum_{n \in \mathcal{M}} R_{\mathcal{P}_n} = \frac{1}{|\mathcal{M}|} \sum_{n \in \mathcal{M}} \left(\prod_{l \in \mathcal{P}_n} R_l\right)$$

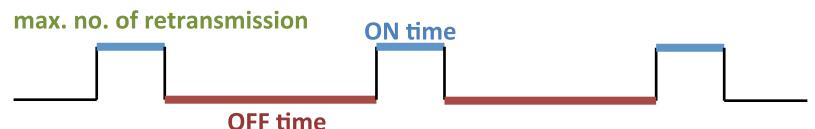
$$Protocol-independent (child-parent link)$$

$$R = \frac{1}{|\mathcal{M}|} \sum_{n \in \mathcal{M}} R_{\mathcal{P}_n} = \frac{1}{|\mathcal{M}|} \sum_{n \in \mathcal{M}} \left(\prod_{l \in \mathcal{P}_n} R_l\right)$$

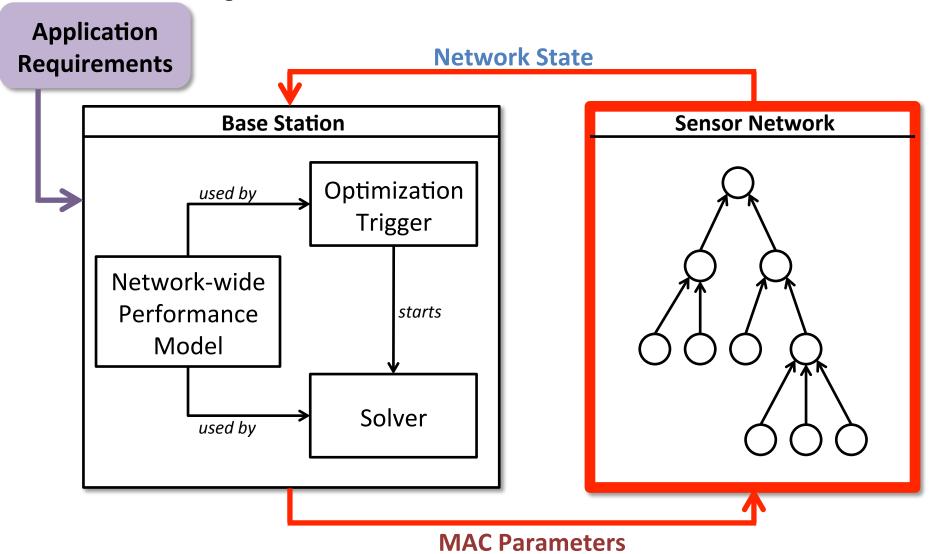
$$R_l = 1 - (1 - p_{s,l})^{N+1}$$

$$X-MAC-specific (single transmission)$$

$$p_{s,l} = \left(1 - (1 - p_l)^{(T_{on} - T_{str})/T_{it}}\right) \cdot p_l^2$$



## pTunes in a Nutshell

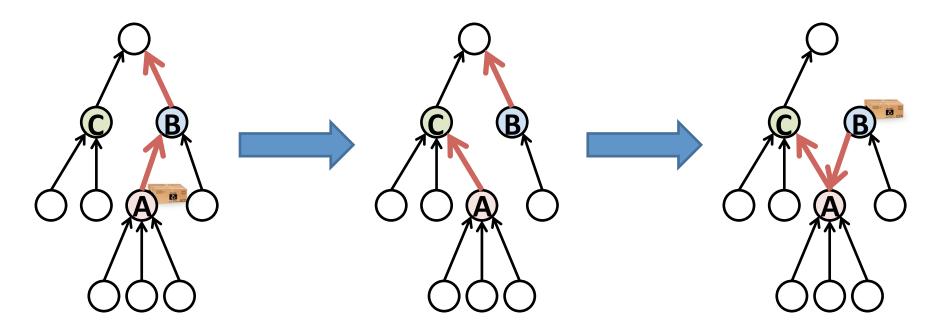


#### **Communication Support to Close the Loop**

- Piggybacking introduces a dependence on the rate and the reliability of application traffic
- Running a dissemination protocol concurrently may degrade the reliability of application traffic

Need to *decouple* network state collection and MAC parameter dissemination from application data traffic

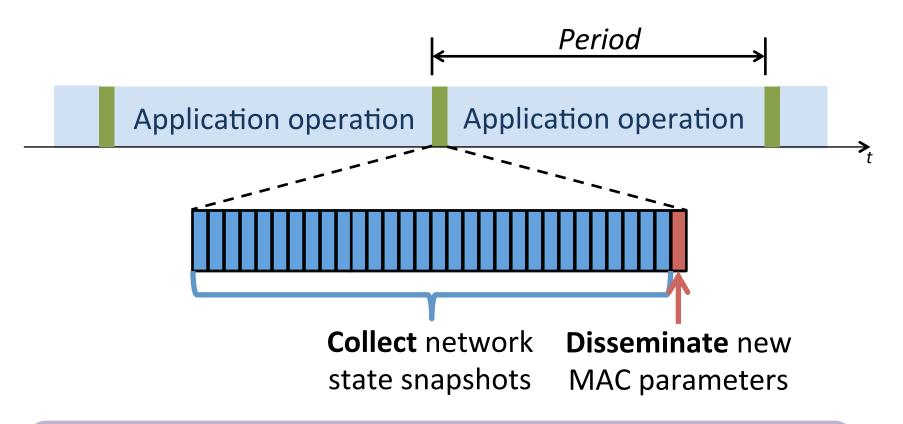
#### The Need for Consistency



A reports B and B reports A as parent, resulting in a loop that never existed for real

Need to collect *consistent* snapshots of network state from all nodes

## Closing the Loop: Our Solution



- Phases consist of *non-overlapping slots*, one for each node
- Each slot corresponds to a distinct Glossy network flood

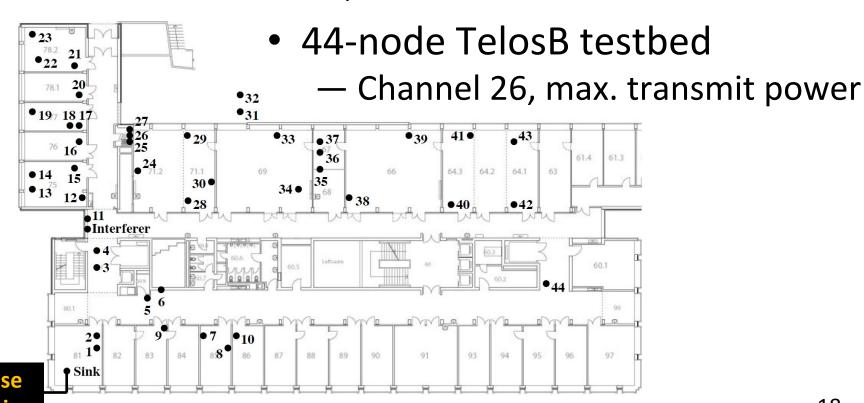
#### **Our Solution Achieves**

- Temporal decoupling from application
- Timeliness
  - Fast collection and dissemination
- Consistency
  - Snapshots taken at all nodes at the same time
  - Simultaneous transition to new MAC parameters
- Energy efficiency (44-node TelosB testbed)

Period	<b>Excess Radio Duty Cycle</b>
1 minute	0.35 %
5 minutes	0.07 %

## **Testbed Evaluation: Setup**

- Implementation
  - Sensor nodes: Contiki, Rime stack, X-MAC, LPP
  - Base station: Java, ECLiPSe



#### **Testbed Evaluation: Methodology**

- Metrics
  - Projected network lifetime: measured in software and computed based on 2000 mAh @ 3V batteries
  - End-to-end reliability: packet sequence numbers
  - End-to-end latency: packet time stamps
- Requirements specification

Maximize: Network lifetime

Subject to: End-to-end reliability greater than 95 %

End-to-end latency below 1 second

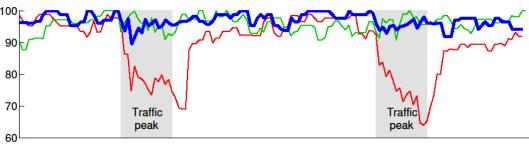
Compare to static MAC parameters determined using pTunes and extensive experiments

#### **Testbed Evaluation: Overview**

- 1. Our X-MAC and LPP models are very accurate
- 2. pTunes provides higher bandwidth against increasing traffic and prevents queue overflows
- 3. pTunes achieves up to three-fold lifetime gains
- 4. Adaptation to traffic fluctuations
- 5. Adaptation to changes in link quality
- 6. Interaction with routing

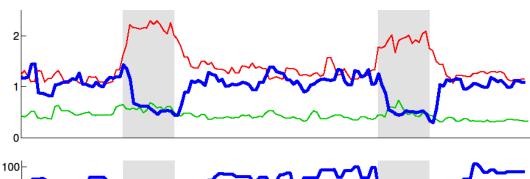
**Adaptation to Traffic Fluctuations** 

End-to-end reliability [%] greater than 95 %



Static 1
Static 2
pTunes

End-to-end latency [sec] below 1 second

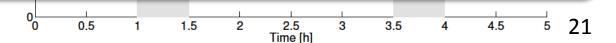


Network lifetime [days]

pTunes satisfies end-to-end requirements at high traffic while extending network lifetime at low traffic

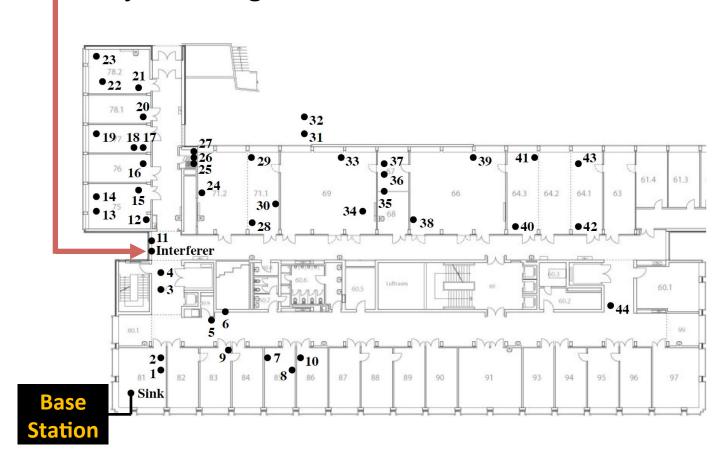
60

20



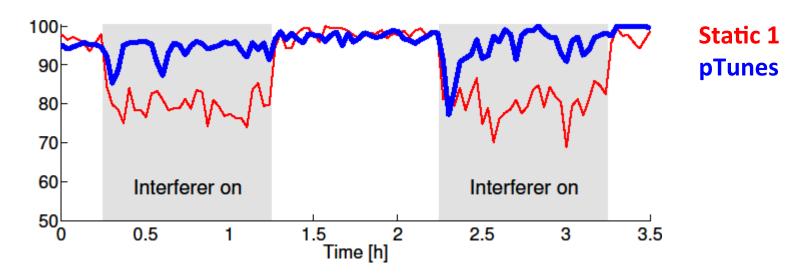
## **Adaptation to Changes in Link Quality**

Interferer jams using modulated carrier: 1 ms ON / 10 ms OFF



## Adaptation to Changes in Link Quality

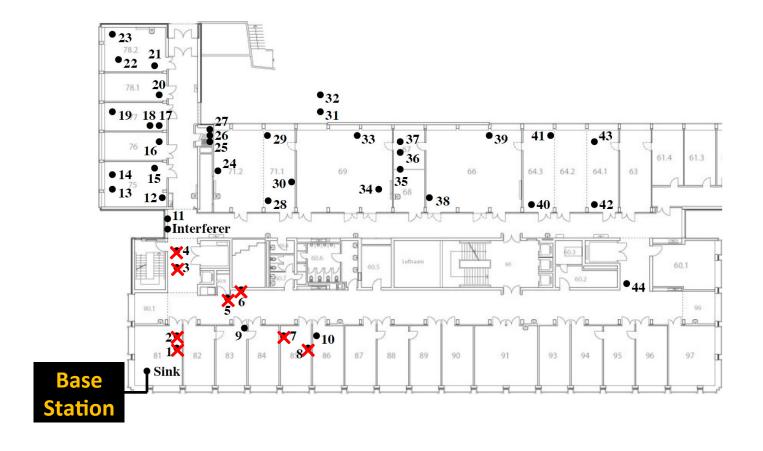
End-to-end reliability [%] greater than 95 %



pTunes reduces packet loss by 80 % during periods of controlled wireless interference

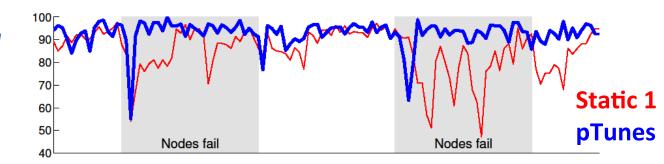
## Interaction with Routing

Turn off 8 nodes × within the sink's neighborhood

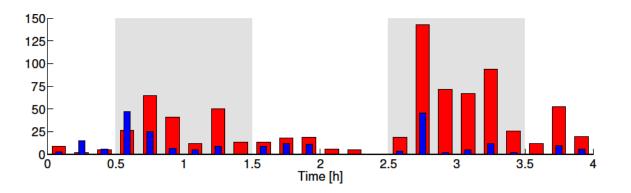


## Interaction with Routing

End-to-end reliability [%] greater than 95 %



Total number of parent switches



pTunes helps the routing protocol quickly recover from node failures, thus reducing packet loss by 70 %

#### **Conclusions**

- pTunes framework for runtime adaptability of existing low-power MAC protocols
- Flexible modeling approach
- Efficient system support to "close the loop"
- Testbed experiments demonstrate that
  - pTunes aids in meeting the requirements of realworld applications as the network state changes
  - pTunes eliminates the need for time-consuming, and yet error-prone, manual MAC configuration









